



Relocate-Vote: Using Sparsity Information to Exploit Ciphertext Side-Channels

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Confidential Computing

Threat model

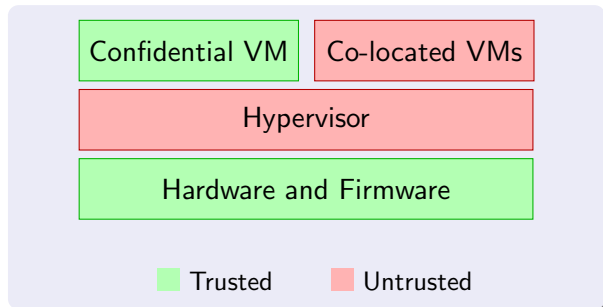
- Untrusted hypervisor
- Trusted hardware and firmware

Cloud service providers (CSPs)



Trusted hardware

AMD SEV-SNP



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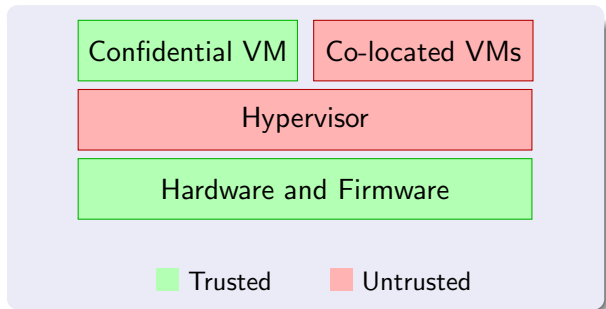
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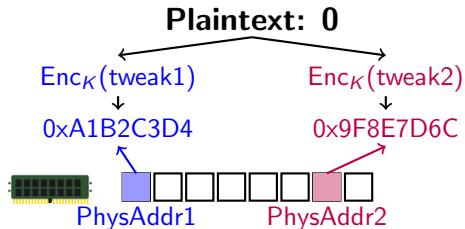
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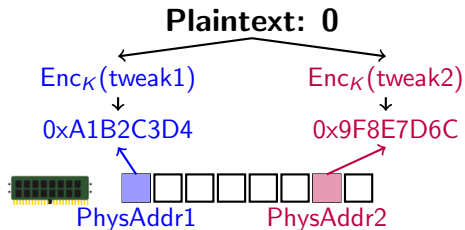
SEV-SNP: Address-dependent Deterministic Memory Encryption

- **Visibility:** Ciphertext is visible to the hypervisor
- Each memory location has unique encryption parameters (tweaks)
 - ▶ Same plaintext \rightarrow different ciphertexts at different locations



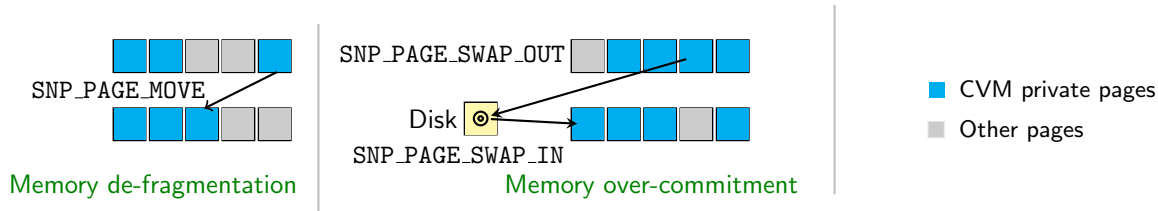
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 - ▶ Same plaintext \rightarrow different ciphertexts at different locations
- **Determinism:** At a fixed location: Same plaintext \rightarrow same ciphertext
 - ▶ Ciphertext side-channels: CIPHERLEAKS [SEC'21], Li et al. [Oakland'22], HyperTheft [CCS'24], CipherSteal [Oakland'25]
 - ★ Only attacks at fixed locations
 - ★ Learned ciphertexts at a location are not useful for other locations

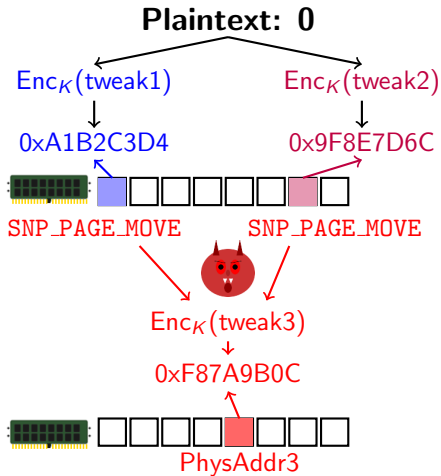


Relocate-Vote: SEV-SNP's Hardware-assisted Page Relocation

- Hardware assists in re-encrypting the relocated pages
- SEV-SNP commands support page relocation
 - ▶ SNP_PAGE_MOVE: Direct relocation for memory de-fragmentation
 - ▶ SNP_PAGE_SWAP: Moving page in/out of disk for memory over-commitment
- **Controlled relocation:** The hypervisor controls the destination page frame

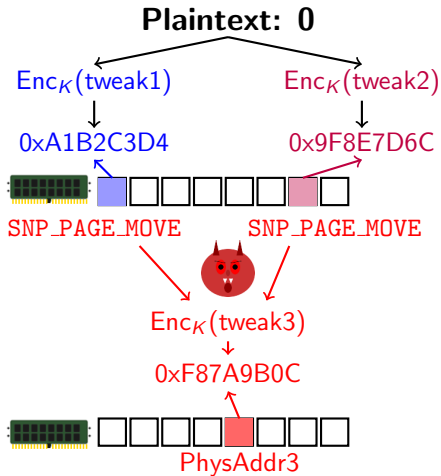


Relocate-Vote: Exploiting Relocation Mechanism



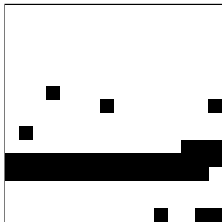
- Plaintext frequency preserved in ciphertext
- Exploit existing values across memory locations
 - ▶ Break tweaked encryption's spatial protection

Relocate-Vote: Exploiting Frequency under Tweaked Encryption



- Infer the ciphertexts of prevalent values (e.g., zero)
 - ▶ Relocate CVM's pages onto the same page frame
 - ▶ Collect the re-encrypted ciphertexts
 - ★ Vote for frequencies
 - ▶ Zero is a prevalent value in CVMs

Relocate-Vote: Exploiting Frequency under Tweaked Encryption



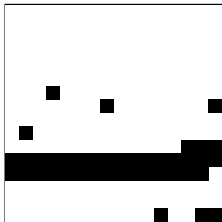
Test a CVM's private page:

□ blocks: zero memory locations

■ blocks: non-zero memory locations

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 - ▶ Examine if ciphertexts match the prevalent ones

Relocate-Vote: Exploiting Frequency under Tweaked Encryption



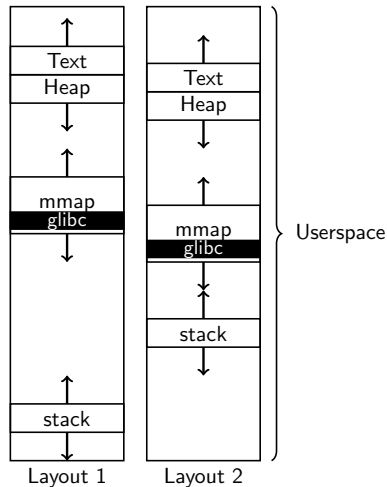
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- Leverage patterns of prevalent and non-prevalent values

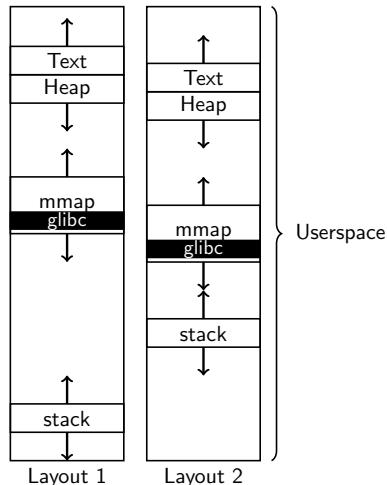
De-randomize ASLR with Relocate-Vote

- ASLR: Address Space Layout Randomization
 - ▶ Randomizes memory layout
 - ▶ Prevents predictable addresses for exploitation, such as `glibc`'s addresses for ROP.

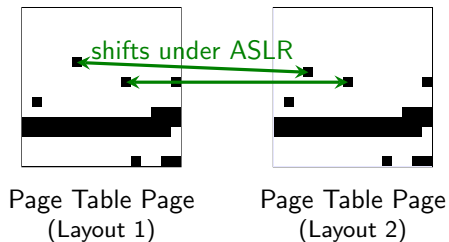


De-randomize ASLR with Relocate-Vote

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 - ▶ Randomizes memory layout
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- Our attack infers `glibc`'s location:
 - ▶ Base address (guest virtual address, GVA)
 - ▶ Queries services provided by the victim application
 - ▶ No direct code execution inside the CVM



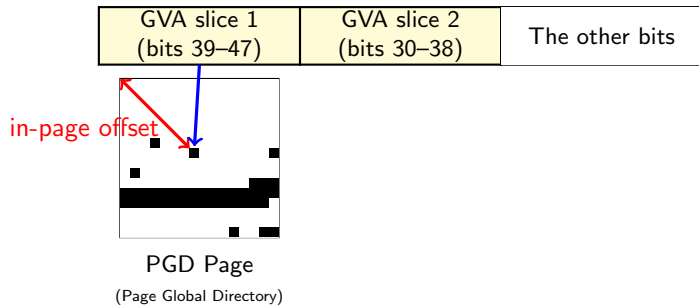
De-randomize ASLR: Sparsity in Page Tables



- Mapped regions (non-zeroed entries) alternate with unmapped regions (zeroed entries)
 - ▶ ASLR shifts the mapped regions → shifts the non-zero blocks

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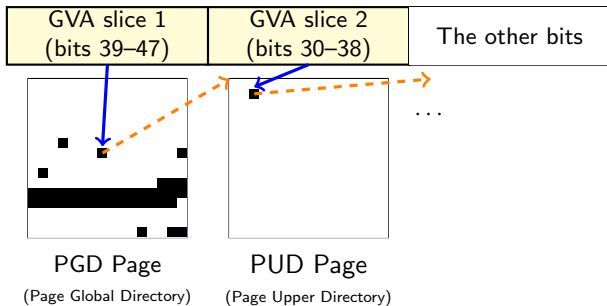
Guest Virtual Address (GVA as the **secret**)



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- GVA slices index page table entries → Offsets of entries reveal GVA slices

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- Address translation walks through page table levels: PGD → PUD → PMD → PTE

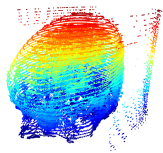
De-randomize ASLR: End-to-end Attack and Results

- Offline phase
 - ▶ Profiles on attacker's own CVM
 - ▶ Prepares the classifiers for identifying page table pages and producing offsets
- Online phase
 - ▶ Triggers the service accessing `glibc` symbols provided by the victim CVM
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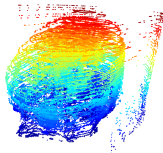
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- **Results:** 8,388,608 possibilities reduced to 35-104 on average
 - ▶ 5 server applications (nginx, apache, memcached, redis, and mysql)
 - ▶ 10 different memory layouts per application

Other Attack Scenarios with Relocate-Vote



Original

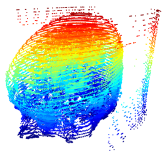


Recovered

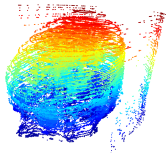
Extract 3D object constructed from CT scanning

- OpenVDB: Library for representing and processing sparse 3D voxel grids
 - ▶ Extracts voxel distributions from the victim's construction and read-only traversal operations

Other Attack Scenarios with Relocate-Vote

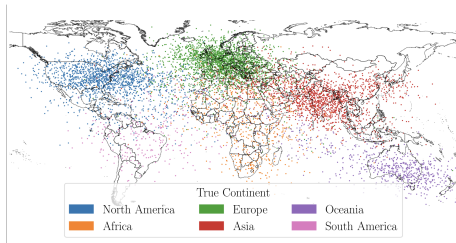


Original



Recovered

Extract 3D object constructed from CT scanning



Predicted coordinates of the places from activation patterns

- OpenVDB: Library for representing and processing sparse 3D voxel grids
 - ▶ Extracts voxel distributions from the victim's construction and read-only traversal operations
- Sparse LLM: LLM with ReLU activations
 - ▶ Decodes geographical information about the processed prompts from activation patterns
- More details in our paper

Mitigation

- Mitigating sparsity leakage at the software level
 - ▶ ↓ Compatibility and performance
- Restricting the hypervisor's relocation ability
 - ▶ **A new guest policy in SEV-SNP:** `PAGE_SWAP_DISABLE`
 - ▶ ↓ Hypervisor's ability of memory resources management
- Enforcing ciphertext-hiding in SEV-SNP instances
 - ▶ ↓ Specific hardware required
 - ▶ ↓ Limited availability of supported SEV-SNP instances in major CSPs

Conclusion

- **Relocate** (page relocation) and **Vote** (frequency analysis)
 - ▶ Leaks ciphertexts of prevalent values
 - ▶ Tests arbitrary memory locations
- Recovered sparsity information in CVMs
 - ▶ ASLR, OpenVDB, sparse LLM
 - ▶ Various operations: lookup, construction, traversal
- Exacerbated implications for ciphertext side channels
 - ▶ Broadens attack scenarios—no ciphertext updates or collisions
 - ▶ Cross-location ciphertext knowledge transfer
- Mitigation is required
 - ▶ Security brief **AMD-SB-3021** by AMD